CS212 – Operating Systems

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Stanford University

Outline

1. Administrivia
2. Substance

CS212 vs. CS112

- **CS212 (previously CS140) is a standalone OS class**
  - Lectures introduce OS topics, similar to CS111
  - Exams test you on material from lecture
  - Programming projects make ideas concrete in an instructional OS
- **CS112 is just the projects from CS212**
  - Only makes sense if you've previously taken CS111
  - Idea: projects in separate quarter from lectures allows more time
  - Feel free to attend any lectures if you want to review a topic (but most will be similar to CS111)
  - A few recommended lectures/sections marked in syllabus
- **In case there are still bugs in program sheets**
  - CS111 or CS212 should fulfill any OS breadth requirement
  - CS112 or CS212 should satisfy significant implementation
  - Ask for exception if something doesn't make sense

Lecture attendance

- **In-person lecture attendance expected of most CS212 students**
- **Exceptions**
  - You are an SCPD student (welcome to attend but not required)
  - Lecture conflicts with another class for which attendance required
  - Occasional one-of conflicts (travel, COVID, sports competitions)
- **Don't just watch the videos if you are an in-person student**
  - Especially don't save all the videos until the night before the exam
- **Lectures will be available by zoom and recorded**
  - When practical, SCPD encouraged to join synchronously via zoom
  - Otherwise, videos will be on panopto

Administrivia

- **Class web page:** `http://cs212.scs.stanford.edu/`
  - All assignments, handouts, lecture notes on-line
- **Textbook: *Operating System Concepts, 8th Edition*, by Silberschatz, Galvin, and Gagne**
  - Out of print and highly optional (weening class from textbook)
- **Goal is to make lecture slides the primary reference**
  - Almost everything I talk about will be on slides
  - PDF slides contain links to further reading about topics
  - Please download slides from class web page
  - Will try to post before lecture for taking notes (but avoid calling out answers if you read them from slides)

Administrivia 2

- **Edstem is the main discussion forum**
- **Staff mailing list:** `cs212-staff@scs.stanford.edu`
  - Please use edstem for any questions others could conceivably have
- **CA split office hours, first round-robin, then individual group**
  - Please ask non-private questions in RR portion
  - Priority for individual group will go to people who attended RR
- **Key dates:**
  - Lectures: MW 1:30pm–2:50pm
  - Section: 6 Fridays, starting this Friday (time, location TBD)
  - Midterm: Monday, February 12, in class (1:30pm-2:50pm)
  - Final: Wednesday, March 20, 3:30pm-6:30pm
  - In-person attendance required for midterm and final (except SCPD)
  - SCPD can use exam monitor, return within 24 hours of exam start
- **Exams open note, but not open book**
  - Bring notes, slides, any printed materials *except* textbook

Course topics

- Threads & Processes
- Concurrency & Synchronization
- Scheduling
- Virtual Memory
- I/O
- Disks, File systems
- Protection & Security
- Virtual machines
- Note: Lectures will often take Unix as an example
  - Most current and future OSes heavily influenced by Unix
  - Won't talk much about Windows

Course goals

- Introduce you to operating system concepts
  - Hard to use a computer without interacting with OS
  - Understanding the OS makes you a more effective programmer
- Cover important systems concepts in general
  - Caching, concurrency, memory management, I/O, protection
- Teach you to deal with larger software systems
  - Programming assignments much larger than many courses
  - Warning: Many people will consider course very hard
  - In past, majority of people report $\geq$15 hours/week
  - We hope it's more manageable with CS111 background and no lectures or exams
- Prepare you to take graduate OS classes (CS240, 240[a-z])

Programming Assignments

- Implement parts of Pintos operating system
  - Built for x86 hardware, you will use hardware emulators
- One setup homework (lab 0) due this Friday
- Four two-week implementation projects:
  - Threads
  - User processes
  - Virtual memory
  - File system
- Lab 1 distributed at end of this week
  - Attend section this Friday for project 1 overview
- Implement projects in groups of up to 3 people
  - CS112/CS212 mixed groups allowed
  - Disclose to partners if you are taking class pass/fail
  - Use "Forming Teams" category on edstem to meet people

Grading

- No incompletes
  - Talk to instructor ASAP if you run into real problems
- Final grades posted March 26
- 50% of CS212 grade based on exams using this quantity:
  $\max\left(\textbf{midterm} > 0\ ?\ \textbf{final} : 0, \frac{1}{2}(\textbf{midterm} + \textbf{final})\right)$
- 50% of CS212 grade, 100% of CS112 grade from projects
  - For each project, 50% of score based on passing test cases
  - Remaining 50% based on design and style
- Most people's projects pass most test cases
  - Please, please, please turn in working code, or **no credit** here
- Means design and style matter a lot
  - Large software systems not just about producing working code
  - Need to produce code other people can understand
  - That's why we have group projects

Style

- Must turn in a design document along with code
  - We supply you with templates for each project's design doc
- CAs will manually inspect code for correctness
  - E.g., must actually implement the design
  - Must handle corner cases (e.g., handle `malloc` failure)
- Will deduct points for error-prone code w/o errors
  - Don't use global variables if automatic ones suffice
  - Don't use deceptive names for variables
- Code must be easy to read
  - Indent code, keep lines and (when possible) functions short
  - Use a uniform coding style (try to match existing code)
  - Put comments on structure members, globals, functions
  - Don't leave in reams of commented-out garbage code

Assignment requirements

- Do not look at other people's solutions to projects
  - We reserve the right to run MOSS on present and past submissions
  - Do not publish your own solutions in violation of the honor code
  - That means using (public) github can get you in big trouble
- You may read but not copy other OSes
  - E.g., Linux, OpenBSD/FreeBSD, etc.
- Cite any code that inspired your code
  - As long as you cite what you used, it's not cheating
  - In worst case, we deduct points if it undermines the assignment
- Projects due 30 minutes before section Fridays
  - Free extension to 5pm if you attend/watch section
- Ask `cs212-staff` for extension if you run into trouble
  - Be sure to tell us: How much have you done? How much is left? When can you finish by?

Outline

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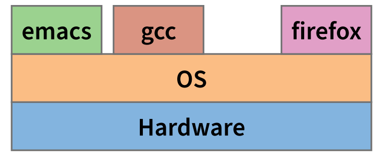
What is an operating system?

- **Layer between applications and hardware**



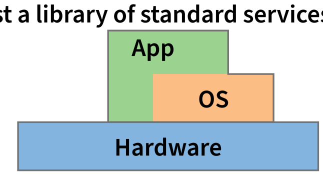
- **Makes hardware useful to the programmer**
- **[Usually] Provides abstractions for applications**
  - Manages and hides details of hardware
  - Accesses hardware through low/level interfaces unavailable to applications
- **[Often] Provides protection**
  - Prevents one process/user from clobbering another

Why study operating systems?

- **Operating systems are a mature field**
  - Most people use a handful of mature OSes
  - Hard to get people to switch operating systems
  - Hard to have impact with a new OS
- **Still open questions in operating systems**
  - Security – Hard to achieve security without a solid foundation
  - Scalability – How to adapt concepts when hardware scales 10× (fast networks, low service times, high core counts, big data…)
- **High-performance servers are an OS issue**
  - Face many of the same issues as OSes, sometimes bypass OS
- **Resource consumption is an OS issue**
  - Battery life, radio spectrum, etc.
- **New “smart” devices need new OSes**

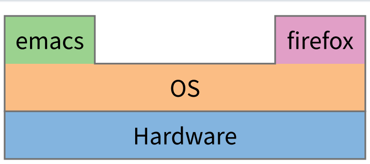
Primitive Operating Systems

- **Just a library of standard services [no protection]**



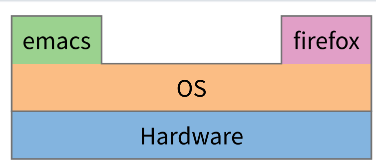
  - Standard interface above hardware-specific drivers, etc.
- **Simplifying assumptions**
  - System runs one program at a time
  - No bad users or programs (often bad assumption)
- **Problem: Poor utilization**
  - …of hardware (e.g., CPU idle while waiting for disk)
  - …of human user (must wait for each program to finish)

Multitasking



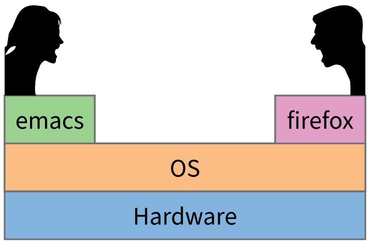
- **Idea: More than one process can be running at once**
  - When one process blocks (waiting for disk, network, user input, etc.) run another process
- **Problem: What can ill-behaved process do?**

Multitasking



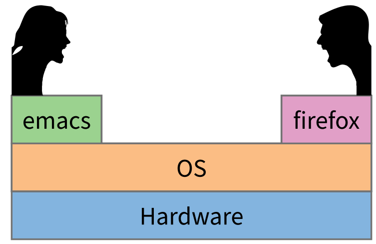
- **Idea: More than one process can be running at once**
  - When one process blocks (waiting for disk, network, user input, etc.) run another process
- **Problem: What can ill-behaved process do?**
  - Go into infinite loop and never relinquish CPU
  - Scribble over other processes’ memory to make them fail
- **OS provides mechanisms to address these problems**
  - *Preemption* – take CPU away from looping process
  - *Memory protection* – protect processes’ memory from one another

Multi-user OSes



- **Many OSes use *protection* to serve distrustful users/apps**
- **Idea: With *N* users, system not *N* times slower**
  - Users' demands for CPU, memory, etc. are bursty
  - Win by giving resources to users who actually need them
- **What can go wrong?**

Multi-user OSes



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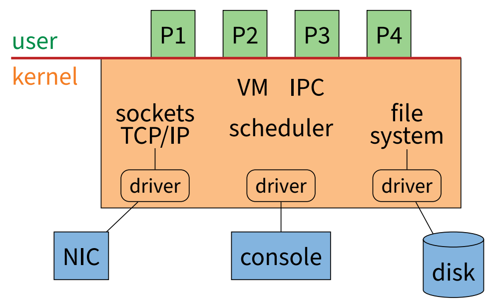
- **Idea: With *N* users, system not *N* times slower**
  - Users' demands for CPU, memory, etc. are bursty
  - Win by giving resources to users who actually need them
- **What can go wrong?**
  - Users are gluttons, use too much CPU, etc. (need policies)
  - Total memory usage greater than machine's RAM (must virtualize)
  - Super-linear slowdown with increasing demand (thrashing)

Protection

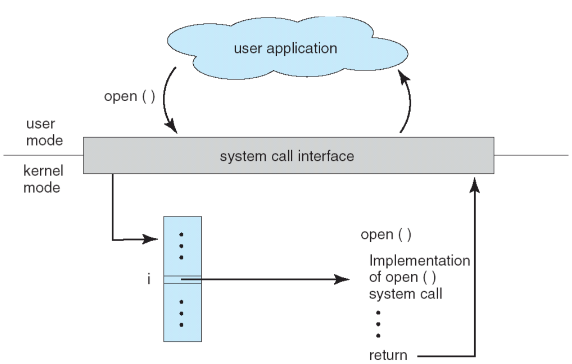
- **Mechanisms that isolate bad programs and people**
- **Pre-emption:**
  - Give application a resource, take it away if needed elsewhere
- **Interposition/mediation:**
  - Place OS between application and "stuff"
  - Track all pieces that application allowed to use (e.g., in table)
  - On every access, look in table to check that access legal
- **Privileged & unprivileged modes in CPUs:**
  - Applications unprivileged (unprivileged *user* mode)
  - OS privileged (privileged supervisor/*kernel* mode)
  - Protection operations can only be done in privileged mode

Typical OS structure



- **Most software runs as user-level processes (P[1-4])**
  - process $\approx$ instance of a program
- **OS *kernel* runs in *privileged* mode (orange)**
  - Creates/deletes processes
  - Provides access to hardware

System calls

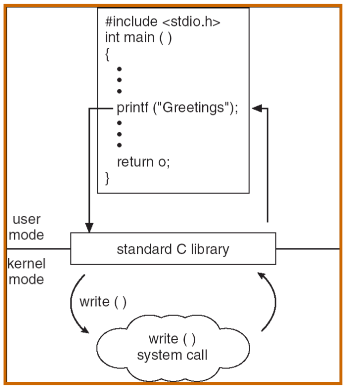


- **Applications can invoke kernel through *system calls***
  - Special instruction transfers control to kernel
  - . . . which dispatches to one of few hundred syscall handlers

System calls (continued)

- **Goal: Do things application can't do in unprivileged mode**
  - Like a library call, but into more privileged kernel code
- **Kernel supplies well-defined *system call* interface**
  - Applications set up syscall arguments and *trap* to kernel
  - Kernel performs operation and returns result
- **Higher-level functions built on syscall interface**
  - `printf`, `scanf`, `fgets`, etc. all user-level code
- **Example: POSIX/UNIX interface**
  - `open`, `close`, `read`, `write`, ...

System call example



- **Standard library implemented in terms of syscalls**
  - *printf* – in libc, has same privileges as application
  - calls *write* – in kernel, which can send bits out serial port

UNIX file system calls

- **Applications "open" files (or devices) by name**
  - I/O happens through open files
- `int open(char *path, int flags, /*int mode*/...);`
  - `flags`: `O_RDONLY`, `O_WRONLY`, `O_RDWR`
  - `O_CREAT`: create the file if non-existent
  - `O_EXCL`: (w. `O_CREAT`) create if file exists already
  - `O_TRUNC`: Truncate the file
  - `O_APPEND`: Start writing from end of file
  - `mode`: final argument with `O_CREAT`
- **Returns file descriptor—used for all I/O to file**

Error returns

- **What if `open` fails? Returns -1 (invalid fd)**
- **Most system calls return -1 on failure**
  - Specific kind of error in global int `errno`
  - In retrospect, bad design decision for threads/modularity
- `#include <sys/errno.h>` **for possible values**
  - 2 = `ENOENT` "No such file or directory"
  - 13 = `EACCES` "Permission Denied"
- `perror` **function prints human-readable message**
  - `perror ("initfile");`
    → "`initfile: No such file or directory`"

Operations on file descriptors

- `int read (int fd, void *buf, int nbytes);`
  - Returns number of bytes read
  - Returns 0 bytes at end of file, or -1 on error
- `int write (int fd, const void *buf, int nbytes);`
  - Returns number of bytes written, -1 on error
- `off_t lseek (int fd, off_t pos, int whence);`
  - `whence`: 0 – start, 1 – current, 2 – end
    - Returns previous file offset, or -1 on error
- `int close (int fd);`

File descriptor numbers

- **File descriptors are inherited by processes**
  - When one process spawns another, same fds by default
- **Descriptors 0, 1, and 2 have special meaning**
  - 0 – "standard input" (`stdin` in ANSI C)
  - 1 – "standard output" (`stdout`, `printf` in ANSI C)
  - 2 – "standard error" (`stderr`, `perror` in ANSI C)
  - Normally all three attached to terminal
- **Example:** `type.c`
  - Prints the contents of a file to `stdout`

`type.c`

```
void
typefile (char *filename)
{
  int fd, nread;
  char buf[1024];

  fd = open (filename, O_RDONLY);
  if (fd == -1) {
    perror (filename);
    return;
  }

  while ((nread = read (fd, buf, sizeof (buf))) > 0)
    write (1, buf, nread);

  close (fd);
}
```

- **Can see system calls using strace utility (ktrace on BSD)**

Protection example: CPU preemption

- **Protection mechanism to prevent monopolizing CPU**
- **E.g., kernel programs timer to interrupt every 10 ms**
  - Must be in supervisor mode to write appropriate I/O registers
  - User code cannot re-program interval timer
- **Kernel sets interrupt to vector back to kernel**
  - Regains control whenever interval timer fires
  - Gives CPU to another process if someone else needs it
  - Note: must be in supervisor mode to set interrupt entry points
  - No way for user code to hijack interrupt handler
- **Result: Cannot monopolize CPU with infinite loop**
  - At worst get $1/N$ of CPU with N CPU-hungry processes

Protection is not security

- **How *can* you monopolize CPU?**

Protection is not security

- **How *can* you monopolize CPU?**
- **Use multiple processes**
- **For many years, could wedge most OSes with**

```
int main() { while(1) fork(); }
```

  - Keeps creating more processes until system out of proc. slots
- **Other techniques: use all memory (`chill` program)**
- **Typically solved with technical/social combination**
  - Technical solution: Limit processes per user
  - Social: Reboot and yell at annoying users
  - Social: Ban harmful apps from play store

Address translation

- **Protect memory of one program from actions of another**
- **Definitions**
  - *Address space*: all memory locations a program can name
  - *Virtual address*: addresses in process' address space
  - *Physical address*: address of real memory
  - *Translation*: map virtual to physical addresses
- **Translation done on every load, store, and instruction fetch**
  - Modern CPUs do this in hardware for speed
- **Idea: If you can't name it, you can't touch it**
  - Ensure one process's translations don't include any other process's memory

More memory protection

- **CPU allows kernel-only virtual addresses**
  - Kernel typically part of all address spaces, e.g., to handle system call in same address space
  - But must ensure apps can't touch kernel memory
- **CPU lets OS disable (invalidate) particular virtual addresses**
  - Catch and halt buggy program that makes wild accesses
  - Make virtual memory seem bigger than physical (e.g., bring a page in from disk only when accessed)
- **CPU enforced read-only virtual addresses useful**
  - E.g., allows sharing of code pages between processes
  - Plus many other optimizations
- **CPU enforced execute disable of VAs**
  - Makes certain code injection attacks harder

Different system contexts

- **At any point, a CPU (core) is in one of several contexts**
- ***User-level* – CPU in user mode running application**
- **Kernel process context – i.e., running kernel code on behalf of a particular process**
  - E.g., performing system call, handling exception (memory fault, numeric exception, etc.)
  - Or executing a kernel-only process (e.g., network file server)
- **Kernel code not associated with a process**
  - Timer interrupt (hardclock)
  - Device interrupt
  - "Softirqs", "Tasklets" (Linux-specific terms)
- **Context switch code – change which process is running**
  - Requires changing the current address space
- **Idle – nothing to do (bzero pages, put CPU in low-power state)**

Transitions between contexts

- **User $\rightarrow$ kernel process context: syscall, page fault, ...**
- **User/process context $\rightarrow$ interrupt handler: hardware**
- **Process context $\rightarrow$ user/context switch: return**
- **Process context $\rightarrow$ context switch: sleep**
- **Context switch $\rightarrow$ user/process context**

Resource allocation & performance

- **Multitasking permits higher resource utilization**
- **Simple example:**
  - Process downloading large file mostly waits for network
  - You play a game while downloading the file
  - Higher CPU utilization than if just downloading
- **Complexity arises with cost of switching**
- **Example: Say disk 1,000 times slower than memory**
  - 1 GiB memory in machine
  - 2 Processes want to run, each use 1 GiB
  - Can switch processes by swapping them out to disk
  - Faster to run one at a time than keep context switching

Useful properties to exploit

- **Skew**
  - 80% of time taken by 20% of code
  - 10% of memory absorbs 90% of references
  - Basis behind cache: place 10% in fast memory, 90% in slow, usually looks like one big fast memory
- **Past predicts future (a.k.a. temporal locality)**
  - What's the best cache entry to replace?
  - If past $\approx$ future, then least-recently-used entry
- **Note conflict between fairness & throughput**
  - Higher throughput (fewer cache misses, etc.) to keep running same process
  - But fairness says should periodically preempt CPU and give it to next process

```c
#include <fcntl.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

void
typefile (char *filename)
{
  int fd, nread;
  char buf[1024];

  fd = open (filename, O_RDONLY);
  if (fd == -1) {
    perror (filename);
    return;
  }

  while ((nread = read (fd, buf, sizeof (buf))) > 0)
    write (1, buf, nread);

  close (fd);
}

int
main (int argc, char **argv)
{
  int argno;
  for (argno = 1; argno < argc; argno++)
    typefile (argv[argno]);
  exit (0);
}
```